

Midterm #1

Please print your name:

No notes, calculators or tools of any kind are permitted. There are 36 points in total. You need to show work to receive full credit.

Good luck!

Problem 1. (6 points) Consider $f(x) = (x+r)(x^2-1)$ where r is some constant. Suppose we want to use Newton's method to calculate the root $x^* = 1$.

- (a) For which values of r is Newton's method guaranteed to converge (at least) quadratically to $x^* = 1$?
- (b) For which values of r does Newton's method converge to $x^* = 1$ faster than quadratically?

Solution. Recall that we showed that, if $f(x^*) = 0$ and $f'(x^*) \neq 0$, then Newton's method (locally) converges to x^* quadratically with rate $\frac{1}{2}|f''(x^*)/f'(x^*)|$.

- (a) Newton's method is guaranteed to converge to 1 provided that $f'(1) \neq 0$.

Since $f(1) = 0$, we have $f'(1) = 0$ if and only if 1 is a repeated root which happens if and only if $r = -1$.

Alternatively, we could compute $f'(x) = (x^2-1) + 2x(x+r)$ so that $f'(1) = 2(r+1)$. Thus $f'(1) = 0$ if and only if $r = -1$.

In either case, we conclude that Newton's method converges (at least) quadratically to $x^* = 1$ if $r \neq -1$.

- (b) Newton's method converges to 1 faster than quadratic if $f'(1) \neq 0$ (i.e. $r \neq -1$) and $f''(1) = 0$.

Continuing the computation from the first part, we calculate $f''(x) = 6x + 2r$.

Thus $f''(1) = 6 + 2r = 0$ if and only if $r = -3$.

Hence, Newton's method converges to 1 faster than quadratic if $r = -3$.

Problem 2. (3 points)

- (a) Give one advantage of the secant method over the regula falsi method.
- (b) Give one advantage of the regula falsi method over the secant method.

Solution.

- (a) An advantage of the secant method is that, if it converges, it typically converges faster than linear. (Also, it does not require an initial interval that is guaranteed to contain a root.)
- (b) An advantage of the regula falsi method is that it is guaranteed to converge.

Problem 3. (6 points) Determine all fixed-points of $f(x) = \frac{x}{x+2}$. For each fixed-point x^* determine whether fixed-point iteration of $f(x)$ converges locally to x^* . If so, determine the exact order of convergence as well as the rate.

Solution. Such values of C are necessarily fixed points of $f(x)$. To find these, we solve $\frac{x}{x+2} = x$.

The resulting fixed points are $x^* = 0$ and $x^* = -1$.

$$f'(x) = \frac{(x+2) - x}{(x+2)^2} = \frac{2}{(x+2)^2}$$

- $f'(0) = \frac{1}{2}$

Since $|f'(0)| < 1$, fixed-point iteration converges locally to 0. The convergence is linear with rate $|f'(0)| = \frac{1}{2}$.

- $f'(-1) = 2$

Since $|f'(-1)| > 1$, fixed-point iteration does not converge locally to -1 .

Problem 4. (3 points) Express $19/6$ in base 2. If necessary, indicate which digits repeat.

Solution. Note that $19/6 = 3 + 1/6$ so that $19/6 = (11.\dots)_2$ with $1/6$ to be accounted for.

- $2 \cdot 1/6 = \boxed{0} + 1/3$

- $2 \cdot 1/3 = \boxed{0} + 2/3$

- $2 \cdot 2/3 = \boxed{1} + 1/3$ and now things repeat...

Hence, $19/6 = (11.001\dots)_2$ and the final two digits 01 repeat: $19/6 = (11.0010101\dots)_2$

Problem 5. (8 points) We wish to compute the root $\sqrt{3}$ of $f(x) = x^2 - 3$ using the bisection method.

- Starting with the interval $[1, 2]$, apply two iterations of bisection. What is the resulting approximation of $\sqrt{3}$?
- After how many iterations can we guarantee that the absolute error is less than 0.001?
- Describe in a few words how the regula falsi method proceeds different from the bisection method.
- Newton's method applied to $x^2 - 3$ is equivalent to fixed-point iteration of which function $g(x)$?

Solution.

(a) Note that $f(1) = -2 < 0$ while $f(2) = 1 > 0$. Hence, $f(x)$ must indeed have a root in the interval $[1, 2]$.

- The midpoint of $[1, 2]$ is $\frac{1+2}{2} = \frac{3}{2}$. Since $f(\frac{3}{2}) = \frac{9}{4} - 3 = -\frac{3}{4} < 0$, a root of $f(x)$ must be in $[\frac{3}{2}, 2]$.
- The midpoint of $[\frac{3}{2}, 2]$ is $\frac{3/2+2}{2} = \frac{7}{4}$. Since $f(\frac{7}{4}) = \frac{49}{16} - 3 = \frac{1}{16} > 0$, a root of $f(x)$ must be in $[\frac{3}{2}, \frac{7}{4}]$.

The best approximation of $\sqrt{3} \approx 1.732$ at this point is the midpoint of the final interval: $\frac{13}{8} \approx 1.625$

- The width of the interval after n steps will be exactly $\ell = \frac{2-1}{2^n} = \frac{1}{2^n}$. Since $\sqrt{3}$ is contained in this interval, the absolute error of approximating it with the midpoint is at most $\ell/2 = \frac{1}{2^{n+1}}$. We need to select n so that $\frac{1}{2^{n+1}} < 10^{-3}$. Knowing that $2^{10} = 1024$ (and $\frac{1}{1024} < \frac{1}{1000}$), we conclude that we need 9 iterations.
- The regula falsi method proceeds like the bisection method. However, instead of using the midpoint $\frac{a+b}{2}$ of the interval $[a, b]$, it uses the root of the secant line of $f(x)$ through $(a, f(a))$ and $(b, f(b))$.
- Newton's method applied to $f(x) = x^2 - 3$ is equivalent to fixed-point iteration of

$$g(x) = x - \frac{f(x)}{f'(x)} = x - \frac{x^2 - 3}{2x} = \frac{x}{2} + \frac{3}{2x}.$$

Problem 6. (2 points)

- Suppose we use the regula falsi method to compute the root of a function $f(x)$. Several iterations result in the intervals $[1, 2]$, $[\frac{7}{6}, 2]$, $[\frac{44}{37}, 2]$, $[\frac{273}{229}, 2]$.

Based on these, our approximation of the root is

- Express -18 in binary using the two's complement representation with 6 bits.

Solution.

- Our approximation of the root is the most recently updated endpoint: $\frac{273}{229}$

(By the way, the intervals were generated for $f(x) = x^2 + 3x - 5$.)

- Since $18 = (010010)_2$, -18 is represented by 101110 (invert all bits, then add 1).

Alternatively, note that $-18 = -2^5 + 14$ and $14 = (1110)_2$ to arrive at the same representation.

Problem 7. (5 points) Suppose we wish to approximate the function $f(x) = 2x \ln(x)$.

- (a) What is the 2nd Taylor polynomial $p_2(x)$ of $f(x)$ at $x = 1$?
- (b) Provide an upper bound for the error of approximating $f(x)$ by $p_2(x)$ on the interval $[1, 2]$.

Solution.

(a) $f'(x) = 2\ln(x) + 2$

$$f''(x) = \frac{2}{x}$$

Hence, the 2nd Taylor polynomial of $f(x)$ at $x = 1$ is

$$p_2(x) = f(1) + f'(1)(x - 1) + \frac{1}{2}f''(1)(x - 1)^2 = 2(x - 1) + (x - 1)^2.$$

Comment. There is typically no reason to expand this out since this approximation is intended to be used for x of the form $1 + \delta$, where δ is small (i.e. for x close to 1).

- (b) Taylor's theorem implies that

$$f(x) - p_2(x) = \frac{f^{(3)}(\xi)}{3!}(x - 1)^3$$

for some ξ between 1 and x .

We compute that $f'''(x) = -\frac{2}{x^2}$. This function is increasing on $(0, \infty)$ and so, in particular, on $[1, 2]$. Therefore, the maximum absolute value on $[1, 2]$ is taken at $x = 1$ or $x = 2$. Since $|f'''(1)| = 2$ and $|f'''(2)| = \frac{1}{2}$, we conclude that $|f'''(\xi)| \leq 2$.

On the other hand, $|(x - 1)^3| \leq 1^3 = 1$ for all $x \in [1, 2]$.

We therefore conclude that the error on $[1, 2]$ is bounded by

$$|f(x) - p_2(x)| = \left| \frac{f^{(3)}(\xi)}{3!}(x - 1)^3 \right| \leq \frac{2}{3!} \cdot 1^3 = \frac{1}{3}.$$

Comment. In this simple case, we can determine the maximal error exactly (without using Taylor's theorem). Since the function $f(x) - p_2(x)$ is decreasing on the interval $[1, 2]$, starting with the value 0, the maximal error must occur at $x = 2$ and is $|4\ln(2) - 3| \approx 0.227$.

Problem 8. (3 points) Represent -2.5 as a single precision floating-point number according to IEEE 754.

Solution. $-2.5 = -1.25 \cdot 2^1 = \underbrace{-1.01}_{\text{binary}} \cdot 2^1$

The exponent 1 gets stored as $1 + 127 = \boxed{1000,0000}$.

Overall, -2.5 is stored as $\boxed{1 \ 1000,0000 \ 0100,0000,0000,0000,0000,0000}$.

(extra scratch paper)