

Birthday paradox

EG Among $n=35$ people, how likely is it that two have the same birthday? ^m
 hashes

$$1 - \underbrace{\left(1 - \frac{1}{365}\right)\left(1 - \frac{2}{365}\right)\left(1 - \frac{3}{365}\right) \dots \left(1 - \frac{34}{365}\right)}_{\text{probability that all birthdays different}} = 81.4\%$$

$n=50 : 97.0\%$

$n=70 : 99.9\%$

\Rightarrow collisions more frequent than one might expect

$365 \rightarrow M$

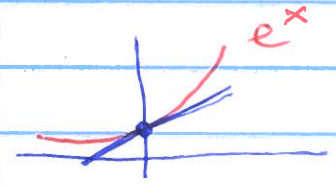
probability of no collision among n selections from M choices =

$e^x e^y = e^{x+y}$

assume $n \ll M$

$$\left(1 - \frac{1}{M}\right)\left(1 - \frac{2}{M}\right) \dots \left(1 - \frac{n-1}{M}\right)$$

$\approx e^{-1/M} \approx e^{-2/M} \approx e^{-(n-1)/M}$



$e^x \approx 1+x$
for small x

$$\approx e^{-(1+2+3+\dots+(n-1))/M}$$

$$= e^{-\frac{(n-1)n}{2M}} \approx e^{-\frac{n^2}{2M}}$$

$$\underbrace{1 + 2 + \dots + (n-1)}_{= \frac{(n-1)n}{2}}$$

if $n \approx \sqrt{M}$

$$\approx e^{-1/2} = 60.7\%$$

(collision likely = 39.3%)

for hashes: output size b bits $\rightarrow M = 2^b$
 among $\sqrt{M} = 2^{b/2}$ hashes
 a collision is likely

\Rightarrow b needs to be twice as large as needed to prevent brute-force

AES-128
key size 128 bit
OK

SHA-2/3 : output size 256+ bits