

Midterm #2

MATH 481/581 — Cryptography
Wednesday, Apr 5

Please print your name:

No notes, calculators or tools of any kind are permitted. There are 36 points in total. You need to show work to receive full credit.

Good luck!

Problem 1. (3+3 points) Bob's public ElGamal key is $(p, g, h) = (19, 10, 6)$.

- (a) Encrypt the message $m = 5$ ("randomly" choose $y = 2$) and send it to Bob.
- (b) Break the cryptosystem and determine Bob's secret key.

Solution.

- (a) The ciphertext is $c = (c_1, c_2)$ with $c_1 = g^y \pmod{p}$ and $c_2 = h^y m \pmod{p}$.

Here, $c_1 = 10^2 \equiv 5 \pmod{19}$ and $c_2 = 6^2 \cdot 5 \equiv -2 \cdot 5 \equiv 9 \pmod{19}$. Hence, the ciphertext is $c = (5, 9)$.

- (b) We need to solve $10^x \equiv 6 \pmod{19}$. Since we haven't learned a better method, we try $x = 1, 2, 3, \dots$ until we find the right one: $10^2 \equiv 5$, $10^3 \equiv 12$, $10^4 \equiv 6 \pmod{19}$. Hence, $x = 4$. \square

Problem 2. (4 points) You are Eve. Alice and Bob select $p = 41$ and $g = 7$ for a Diffie–Hellman key exchange. Alice sends 15 to Bob, and Bob sends 38 to Alice. What is their shared secret?

Solution. Let's crack Alice's secret y (you can also attack Bob; his is $x = 5$).

For that, we need to find y such that $7^y = 15 \pmod{41}$.

We try all possibilities: $7^2 \equiv 8$, $7^3 \equiv 15 \pmod{41}$

Hence, Alice's secret is $y = 3$. Since $38^3 \equiv (-3)^3 \equiv -27 \equiv 14 \pmod{41}$, the shared secret is 14. \square

Problem 3. (3+3 points) Bob's public RSA key is $N = 33$, $e = 13$.

- (a) Encrypt the message $m = 5$ and send it to Bob.
- (b) Determine Bob's secret private key d .

Solution.

- (a) The ciphertext is $c = m^e \pmod{N}$. Here, $c \equiv 5^{13} \pmod{33}$.

$5^2 = 25 \equiv -8$, $5^4 \equiv 64 \equiv -2$, $5^8 \equiv 4 \pmod{33}$. Hence, $5^{13} = 5^8 \cdot 5^4 \cdot 5 \equiv 4 \cdot (-2) \cdot 5 \equiv 26 \pmod{33}$. Hence, $c = 26$.

- (b) $N = 3 \cdot 11$, so that $\phi(N) = 2 \cdot 10 = 20$.

To find d , we compute $e^{-1} \pmod{20}$ using the extended Euclidean algorithm:

$$\begin{aligned} \boxed{20} &= 1 \cdot \boxed{13} + 7 \\ \boxed{13} &= 2 \cdot \boxed{7} - 1 \end{aligned}$$

Backtracking through this, we find that Bézout's identity takes the form

$$1 = 2 \cdot \boxed{7} - \boxed{13} = 2 \cdot (\boxed{20} - 1 \cdot \boxed{13}) - \boxed{13} = 2 \cdot \boxed{20} - 3 \cdot \boxed{13}.$$

Hence, $13^{-1} \equiv -3 \equiv 17 \pmod{20}$ and, so, $d = 17$.

Comment. Bob's choice of $e = 13$ is actually functionally equivalent to $e = 3$ (for instance, $5^3 \equiv 26 \pmod{33}$). Similarly, d can be obtained as $e^{-1} \pmod{10}$. Can you explain these claims? \square

Problem 4. (1+3 points) Consider the finite field $\text{GF}(2^3)$ constructed using $x^3 + x + 1$.

- (a) Multiply $x^2 + 1$ and x in $\text{GF}(2^3)$.
- (b) Determine the inverse of $x^2 + x$ in $\text{GF}(2^3)$.

Solution.

- (a) $(x^2 + 1)x = x^3 + x = 1$ in $\text{GF}(2^3)$.

- (b) In general, we use the extended Euclidean algorithm and reduce modulo 2 at each step. Here, we are lucky and are actually done after a single polynomial division:

$$\boxed{x^3 + x + 1} \equiv (x + 1) \cdot \boxed{x^2 + x} + 1$$

Hence, $(x^2 + x)^{-1} = x + 1$ in $\text{GF}(2^3)$. \square

Problem 5. (16 points) Fill in the blanks.

- (a) If Bob's public RSA key is (N, e) and Bob's secret key is d , then how does Bob decrypt the ciphertext c ?

The plaintext is $m =$

- (b) For his public RSA key, which of p, q and e must Bob choose randomly?

- (c) For his public ElGamal key, which of p, g and x must Bob choose randomly?

- (d) Despite its flaws, in which scenario is it fine to use the Fermat primality test?

- (e) DES has a block size of bits, a key size of bits and consists of rounds.

- (f) To store an S-box in DES as a lookup table, we need bytes.

- (g) Suppose we are using 3DES with key $k = (k_1, k_2, k_3)$, where each k_i is an independent DES key.

Then m is encrypted to $c =$

The effective key size is

bits.

- (h) AES-128 has a block size of bits, a key size of bits and consists of rounds.

- (i) The four layers of AES are

- (j) For his public ElGamal key, Bob selected $p = 101$. He has choices for g .

- (k) For his public RSA key, Bob selected $N = 77$. The smallest choice for e with $e \geq 2$ is

- (l) The primitive roots modulo 5 are

Solution.

- (a) The plaintext is $m = c^d \pmod{N}$.
- (b) p and q must be chosen randomly.
- (c) x must be chosen randomly.
- (d) When testing a huge random number for primality.
- (e) DES has a block size of 64 bits, a key size of 56 bits and consists of 16 rounds.
- (f) Recall that the S-boxes (there is eight different ones) are lookup tables. For each 6 bit input (meaning there is a total of 2^6 possible inputs), they specify 4 bits of output.
To store one S-box, we therefore need to list $2^6 \cdot 4 = 256$ bits, or 32 bytes.
- (g) m is encrypted to $c = E_{k_3}(D_{k_2}(E_{k_1}(m)))$.
The effective key size is 112 bits (because of the meet-in-the-middle attack).
- (h) AES-128 has a block size of 128 bits, a key size of 128 bits and consists of 10 rounds.
- (i) The four layers of AES are: ByteSub, ShiftRow, MixCol, AddRoundKey.
- (j) He has $\phi(100) = 40$ choices for g .
- (k) Since $\phi(77) = 60$, the smallest choice for e with $e \geq 2$ is 7.
- (l) The primitive roots modulo 5 are 2, 3. □

(extra scratch paper)